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CS355: Cryptography

Lecture 3: Vigenere cipher.

Towards polyalphabetic substitution ciphers

- ▶ **Main weaknesses of monoalphabetic substitution ciphers**
 - ▶ each letter in the ciphertext corresponds to only one letter in the plaintext letter
- ▶ **Idea for a stronger cipher (1460' s by Alberti)**
 - ▶ use more than one cipher alphabet, and switch between them when encrypting different letters
- ▶ **Giovani Battista Bellaso published it in 1553**
- ▶ **Developed into a practical cipher by Blaise de Vigenère and published in 1586**

Vigenère cipher

Definition:

Given m , a positive integer, $P = C = (\mathbb{Z}_{26})^n$, and $K = (k_1, k_2, \dots, k_m)$ a key, we define:

Encryption:

$$e_k(p_1, p_2 \dots p_m) = (p_1 + k_1, p_2 + k_2 \dots p_m + k_m) \pmod{26}$$

Decryption:

$$d_k(c_1, c_2 \dots c_m) = (c_1 - k_1, c_2 - k_2 \dots c_m - k_m) \pmod{26}$$

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25

Example:

Plaintext: C R Y P T O G R A P H Y

Key: L U C K L U C K L U C K

Ciphertext: N L A Z E I I B L J J I

Security of Vigenere cipher

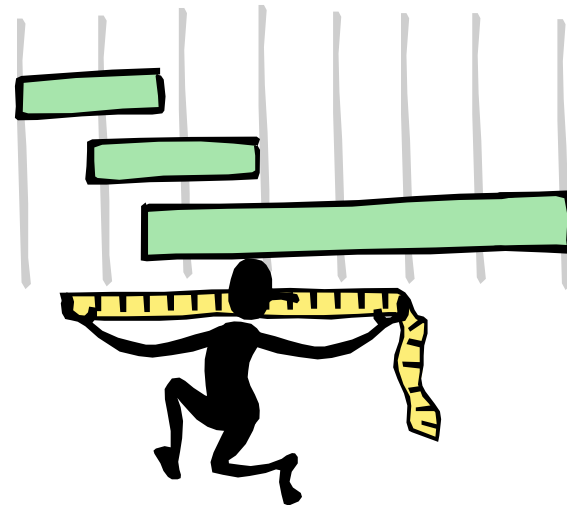
- ▶ Vigenere **masks the frequency** with which a character appears in a language: one letter in the ciphertext corresponds to multiple letters in the plaintext. Makes the **use of frequency analysis more difficult**
- ▶ Any message encrypted by a Vigenere cipher is a collection of as **many shift ciphers** as there are letters in the key

Vigenere cipher cryptanalysis

- ▶ Find the **length of the key**
- ▶ **Divide** the message into that many shift cipher encryptions
- ▶ **Use frequency analysis** to solve the resulting shift ciphers
 - ▶ how?

How to find the key length?

- ▶ For Vigenere, as the length of the keyword increases, the letter frequency shows less English-like characteristics and becomes more random
- ▶ Two methods to find the key length:
 - ▶ Kasisky test
 - ▶ Index of coincidence (Friedman)



History of breaking Vigenere

- ▶ 1596 - Cipher was published by Vigenere
- ▶ 1854 - It is believed the Charles Babbage knew how to break it in 1854, but he did not published the results
- ▶ 1863 - Kasiski showed the Kasiski examination that showed how to break Vigenere
- ▶ 1920 - Friedman published ``The index of coincidence and its applications to cryptography''

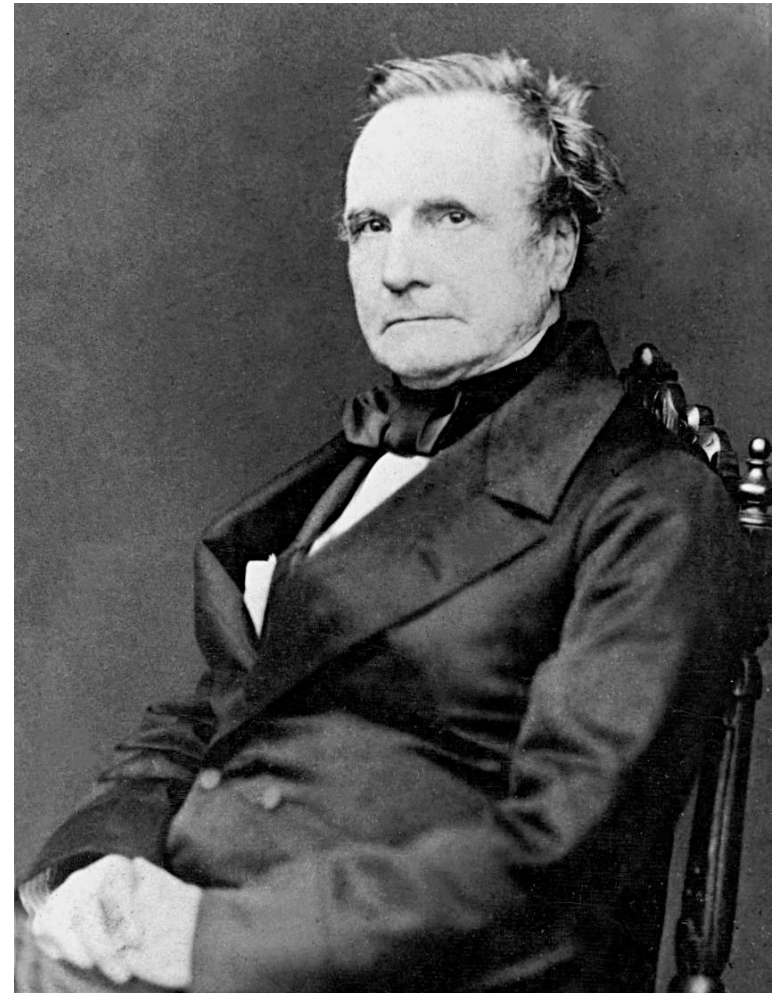
Friedrich Wilhelm Kasiski (1805 – 1881)

- ▶ German infantry officer, cryptographer and archeologist.



Charles Babbage (1791 – 1871)

- ▶ English mathematician, philosopher, inventor and mechanical engineer who originated the concept of a programmable computer.
- ▶ Considered a "father of the computer", he invented the first mechanical computer that eventually led to more complex designs.



William Frederick Friedman (1891 – 1969)

- ▶ US Army cryptographer who ran the research division of the Army's Signals Intelligence Service (SIS) in the 1930s, and parts of its follow-on services into the 1950s.
- ▶ In 1940, people from his group, led by Frank Rowlett broke Japan's PURPLE cipher machine



Kasisky test

- ▶ Note: two identical segments of plaintext, will be encrypted to the same ciphertext, if they occur in the text at the distance Δ , ($\Delta \equiv 0 \pmod{m}$), m is the key length)
- ▶ Algorithm:
 - ▶ Search for pairs of identical segments of length at least 3
 - ▶ Record distances between the two segments: $\Delta_1, \Delta_2, \dots$
 - ▶ m divides $\gcd(\Delta_1, \Delta_2, \dots)$

Example of the Kasisky test

Key K I N G K I N G K I N G K I N G K I N G K
 I N G

PT t h e s u n a n d t h e m a n i n t h e m
 o o n

CT D P R Y E V N T N **B U K** W I A O X **B U K** W
 W B T

Index of coincidence (Friedman)

Informally: Measures the probability that two random elements of the n-letters string x are identical.

Definition:

Suppose $x = x_1x_2\dots x_n$ is a string of n alphabetic characters. Then $I_c(x)$, the index of coincidence is:

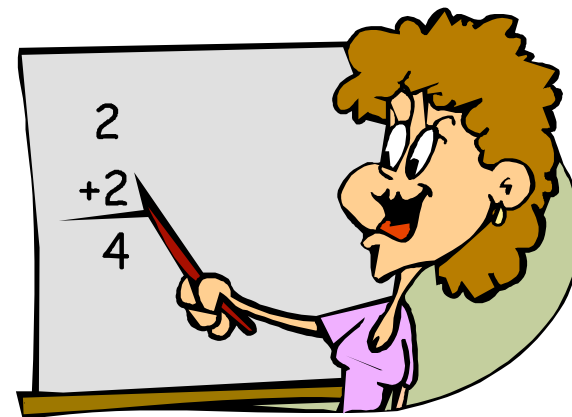
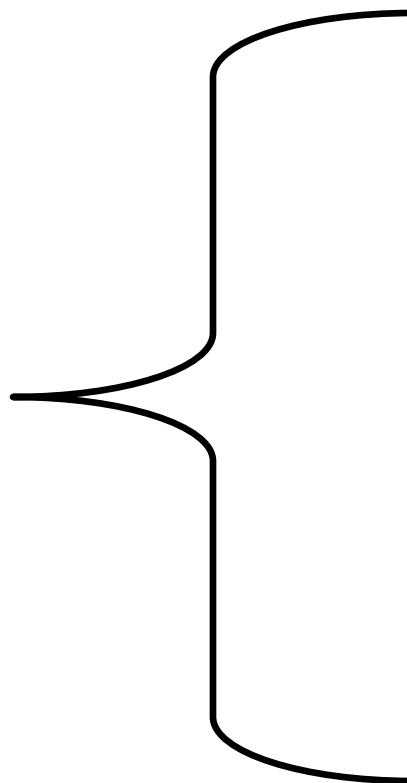
$$I_c(x) = P(x_i = x_j)$$

Index of coincidence (cont.)

- Reminder: binomial coefficient $\binom{n}{k} = \frac{n!}{k!(n-k)!}$
- Consider the plaintext x , and c_0, c_1, \dots, c_{25} are the number of occurrences with which A, B, ... Z appear in x and p_0, p_1, \dots, p_{25} are the probabilities with which A, B, ... Z appear in x .
- We want to compute

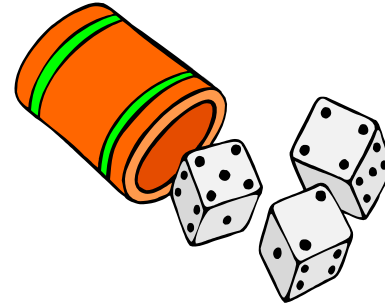
$$I_c(x) = P(x_i = x_j)$$

Begin math



Elements of probability theory

A random experiment has an unpredictable outcome.



Definition

The **sample space (S)** of a random phenomenon is the **set of all outcomes** for a given experiment.

Definition

The **event (E)** is a **subset of a sample space**, an event is any collection of outcomes.

Basic axioms of probability

If E is an event, $Pr(E)$ is the probability that event E occurs then

- (a) $0 \leq Pr(A) \leq 1$ for any set A in S
- (b) $Pr(S) = 1$, where S is the sample space.
- (c) If E_1, E_2, \dots, E_n is a sequence of mutually exclusive events, that is $E_i \cap E_j = \emptyset$, for all $i \neq j$ then:

$$Pr(E_1 \cup E_2 \cup \dots \cup E_n) = \sum_{i=1}^n Pr(E_i)$$

More probabilities

If E is an event and $Pr(E)$ is the probability that the event E occurs then

- ▶ $Pr(\hat{E}) = 1 - Pr(E)$ where \hat{E} is the complimentary event of E
- ▶ If outcomes in S are equally like, then
 $Pr(E) = |E| / |S|$ (where $| |$ denotes the cardinality of the set)

Example

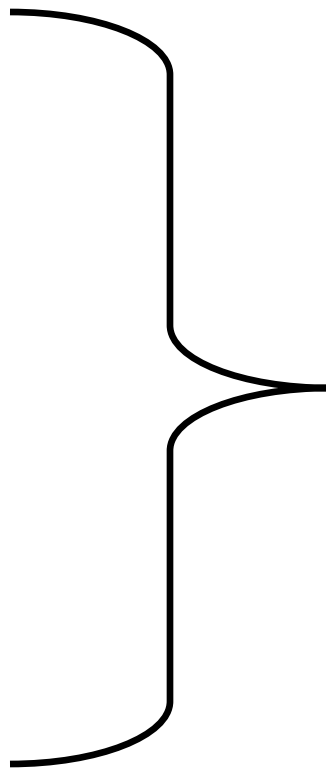
Random throw of a pair of dice.
What is the probability that the sum is 3?

Solution: Each dice can take six different values $\{1,2,3,4,5,6\}$. The number of possible events (value of the pair of dice) is 36, therefore each event occurs with probability $1/36$.

Examine the sum: $3 = 1+2 = 2+1$
The probability that the sum is 3 is $2/36$.

What is the probability that the sum is 11? How about 5?

End math




Index of coincidence (cont.)

- Reminder: binomial coefficient $\binom{n}{k} = \frac{n!}{k!(n-k)!}$
- Consider the plaintext x , and c_0, c_1, \dots, c_{25} are the number of occurrences with which A, B, ... Z appear in x and p_0, p_1, \dots, p_{25} are the probabilities with which A, B, ... Z appear in x .
- We want to compute .

$$I_c(x) = P(x_i = x_j)$$

Index of coincidence (cont.)

- We can choose two elements out of the string of size n in $\binom{n}{2}$ ways
- For each i , there are $\binom{c_i}{2}$ ways of choosing the elements to be i

$$I_C(x) = \frac{\sum_{i=0}^s \binom{c_i}{2}}{\binom{n}{2}} = \frac{\sum_{i=0}^s c_i(c_i - 1)}{n(n-1)} \approx \frac{\sum_{i=0}^s c_i^2}{n^2} = \sum_{i=0}^s p_i^2$$


THIS IS AN APPROXIMATION IF n IS VERY BIG

Example: IC of a string

- ▶ Consider the text **THE INDEX OF COINCIDENCE**

$$I_C(x) = \frac{\sum_{i=0}^s c_i(c_i - 1)}{n(n - 1)}$$

- ▶ There are 21 characters, so $N = 21, S = 25$

$$I_c = (3*2 + 2*1 + 4*3 + 1*0 + 1*0 + 3*2 + 3*2 + 2*1 + 1*0 + 1*0) / 21*20 = \\ \mathbf{34/420 = 0.0809}$$

Example: IC of a language

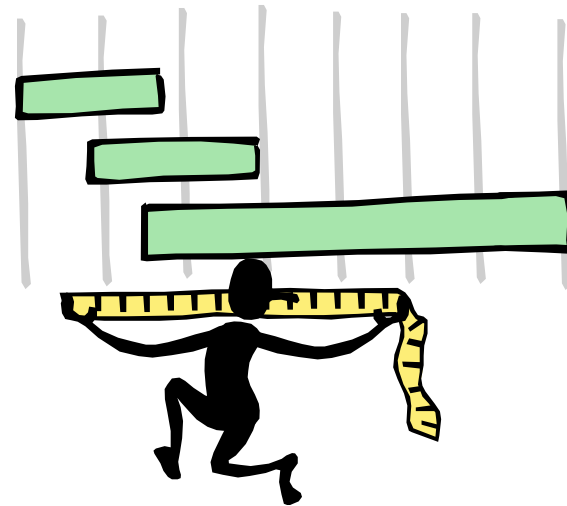
- ▶ For English, $S = 25$ and p_i can be estimated

Letter	p_i	Letter	p_i	Letter	p_i	Letter	p_i
A	.082	H	.061	O	.075	V	.010
B	.015	I	.070	P	.019	W	.023
C	.028	J	.002	Q	.001	X	.001
D	.043	K	.008	R	.060	Y	.020
E	.127	L	.040	S	.063	Z	.001
F	.022	M	.024	T	.091		
G	.020	N	.067	U	.028		

$$I_c(x) = \sum_{i=0}^{i=25} p_i^2 = 0.065$$

Find the key length

- ▶ For Vigenere, as the length of the keyword increases, the letter frequency shows less English-like characteristics and becomes more random.
- ▶ Two methods to find the key length:
 - length:
 - ▶ Kasisky test
 - ▶ Index of coincidence (Friedman)



Finding the key length

$q = q_1 q_2 \dots q_n$, m is the key length

$$\begin{array}{cccc}
 \left[\begin{array}{cccc}
 q_1 & q_{m+1} & \dots & q_{n-m+1} \\
 q_2 & q_{m+2} & \dots & q_{n-m+2} \\
 \dots & \dots & \dots & \dots \\
 q_m & q_{2m} & \dots & q_n
 \end{array} \right] & & \begin{array}{l}
 y_1 \\
 y_2 \\
 \dots \\
 y_m
 \end{array}
 \end{array}$$

Guessing the key length

- ▶ If m is the key length, then the text ``looks like'' **English** text

$$I_c(y_i) \approx \sum_{i=0}^{i=25} p_i^2 = 0.065 \quad \forall 1 \leq i \leq m$$

- ▶ If m is not the key length, the text ``looks like'' **random** text and:

$$I_c \approx \sum_{i=0}^{i=25} \left(\frac{1}{26}\right)^2 = 26 \times \frac{1}{26^2} = \frac{1}{26} = 0.038$$

Finding the key, once key length known

- ▶ Consider vectors y_i , and look for the most frequent letter
- ▶ Check if mapping that letter to e will not result in unlikely mapping for other letters
- ▶ If that's not the case, look at the shift of the mapping, that represents the letter of the key
- ▶ Repeat for each vector

Kasisky example

- ▶ Suppose that a Kasiski analysis of the ciphertext from a Vigenere cipher identifies these seven pairs of repeated sequences of ciphertext letters.

First occurence	8	20	38	48	59	72
Second occurence	32	64	110	104	163	132

- ▶ What can you say about the length of the key used to encrypt the message?

Index of coincidence example

- ▶ A ciphertext of 100 letters was intercepted. The frequency distribution of letters of the alphabet in this ciphertext is as follows:

A 2; B 10; C 2; D 5; E 3; F 8; G 1; H 2; I 2;
J 5; K 1; L 1; M 3; N 2; O 10; P 1; Q 8; R 1;
S 8; T 5; U 2; V 1; W 3; X 5; Y 1; Z 8.

- ▶ What is the index of coincidence of this ciphertext?

Index of coincidence example

- ▶ Suppose there is a language that has only three letters: a,b,c.
 - ▶ frequency of letter a is 0.5
 - ▶ frequency of letter b is 0.3
 - ▶ frequency of letter c is 0.2
- ▶ What is the index of coincidence of the language?

Vigenere example

- ▶ EBBBLCKSYMMKTHPDLSPWLCKVJCKDSYMVLCK
- ▶ 5 21 33
- ▶ Key length divides $\gcd(21-5, 33-21) = \gcd(16, 12) = 4$
- ▶ EBBB
- ▶ LCKS
- ▶ YMMK
- ▶ THPD
- ▶ LSPW
- ▶ LCKV
- ▶ JCKD
- ▶ SYMV
- ▶ LCK

Vigenere challenge

- ▶ KVAESZXYFQZGGEVPQVRISOHVZTXYXOCG
XQVDRKALKIEKUUCAKXKOGKSSSMKXUU
GMTEIJSSPWBGVFBZREWZWVSTVSQZML
NGSURSFYCIINIRGGNUVTLGUWLEEEAYXQC
GLFJGBXLTRTXDXLQVPQKUEXASJFIWDXDLF
TFKHJRLVVULVYHUELFJGBXLXYFVVPLFYFJ
EOGRPRXPRCRLQVPQIEZKPRFUQEMGIXB
- ▶ 10 bonus points, submit all your code you used to solve it
- ▶ Deadline is Sept. 2

Take home lessons

- ▶ Vigenère cipher is vulnerable: once the key length is found, a cryptanalyst can apply frequency analysis.

